





Applied Cryptography CMPS 297AD/396AI Fall 2025

Part 1: Provable Security 1.1: Introduction

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# Defining cryptography

#### What is Cryptography?

"The science of enabling secure and private computation, communication, verification, and delegation in the presence of untrusted parties, adversarial behavior, and mutually distrustful participants."



Source: Serious Cryptography, 2nd Edition

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## Cryptography is everywhere

- Banking
- Buying stuff from the store
- Any digital payment system
- Messaging (WhatsApp, Signal, iMessage, Telegram)
- Voice calls
- Government and military systems
- SSH
- VPN access
- Visiting most websites (HTTPS)

- Disk encryption
- Cloud storage
- Video conferencing
- Unlocking your (newer) car
- Identity card systems
- Ticketing systems
- DRM solutions
- Private contact discovery
- Cryptocurrencies
- That iPhotos feature that detects similar photos

#### How it's made



Fischer et al., The Challenges of Bringing Cryptography from Research Papers to Products: Results from an Interview Study with Experts, USENIX Security 2024

Applied Cryptography - American University of Beirut

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## Cryptographic building blocks

#### Components

- Cryptography manifests as a set of primitives, from which we build protocols intended to accomplish well-defined security goals.
- Primitives: AES, RSA, SHA-2, DH...
- **Protocols**: TLS, Signal, SSH, FileVault 2, BitLocker...

#### Examples

- AES: Symmetric encryption
  - Enc(k, m) = c, Dec(k, c) = m.
- SHA-2: Hash function
  - H(m) = h.
- Diffie-Hellman: Public key agreement
  - Allows two parties to agree on a secret key *k*.

# Cryptographic building blocks

#### Security goals

- **Confidentiality**: Data exchanged between Client and Server is only known to those parties.
- Authentication: If Server receives data from Client, then Client sent it to Server.
- Integrity: If Server modifies data owned by Client, Client can find out.

#### Examples

- **Confidentiality**: When you send a private message on Signal, only you and the recipient can read the content.
- Authentication: When you receive an email from your boss, you can verify it actually came from them.
- Integrity: Your computer can verify that software update downloads haven't been tampered with during transmission.

### Security goals: more examples

- **TLS (HTTPS)** ensures that data exchanged between the client and the server is confidential and that parties are authenticated.
  - Allows you to log into gmail.com without your ISP learning your password.
- FileVault 2 ensures data confidentiality and integrity on your MacBook.
  - Prevents thieves from accessing your data if your MacBook is stolen.
- Signal implements post-compromise security, an advanced security goal.
  - Allows a conversation to "heal" in the event of a temporary key compromise.
  - More on that later in the course.

## Why bother?

- Can't we just use access control?
- Strictly speaking, usernames and passwords can be implemented without cryptography...
- Server checks if the password matches, or if the IP address matches, etc. before granting access.
- What's so bad about that?

#### The Problem with Traditional Access Control

- Requires trusting the server completely
- No protection during transmission
- No way to verify integrity
- No way to establish trust between strangers

## The magic of cryptography

#### Cryptography lets us achieve what seems impossible

- Secure communication over insecure channels
- Verification without revealing secrets
- Proof of computation without redoing it

### Hard problems

- Cryptography is largely about equating the security of a system to the difficulty of solving a math problem that is thought to be computationally very expensive.
- With cryptography, we get security systems that we can literally mathematically prove as secure (under assumptions).
- Also, this allows for actual magic.
  - Alice and Bob meet for the first time in the same room as you.
  - You are listening to everything they are saying.
  - Can they exchange a secret without you learning it?

## Time for actual magic



### No known feasible computation

- The discrete logarithm problem:
  - Given a finite cyclic group G, a generator  $g \in G$ , and an element  $h \in G$ , find the integer x such that  $g^x = h$
- In more concrete terms:
  - Let p be a large prime and let g be a generator of the multiplicative group Z<sup>\*</sup><sub>p</sub> (all nonzero integers modulo p).
  - Given:
    - $g \in \mathbb{Z}_p^*, h \in \mathbb{Z}_p^*$
    - Find  $x \in \{0, 1, \dots, p-2\}$  such that  $g^x \equiv h \pmod{p}$
  - This problem is believed to be computationally hard when p is large and g is a primitive root modulo p.
    - "Believed to be" = we don't know of any way to do it that doesn't take forever, unless we have a strong, stable quantum computer (Shor's algorithm)

### Hard problems

#### **Asymmetric Primitives**

- Diffie-Hellman, RSA, ML-KEM, etc.
- "Asymmetric" because there is a "public key" and a "private key" for each party.
- Algebraic, assume the hardness of mathematical problems (as seen just now.)

#### Symmetric Primitives

- AES, SHA-2, ChaCha20, HMAC...
- "Symmetric" because there is one secret key.
- Not algebraic but unstructured, but on their understood resistance to *n* years of cryptanalysis.
- Can act as substitutes for assumptions in security proofs!
  - Example: hash function assumed to be a "random oracle"

### Kerckhoff's principle

- "A cryptosystem should be secure even if everything about the system, except the key, is public knowledge." Auguste Kerckhoffs, 1883
- Why it matters:
  - No "security through obscurity"
  - The key is the only secret: the rest can be audited, tested, trusted
  - Encourages open standards and peer review
  - If your system's security depends on nobody knowing how it works, it's not secure.

## Symmetric primitive example: hash functions

#### **Hash Function Properties**

- Takes input of any size[<+->]
- Produces output of fixed size
- Is deterministic (same input → same output)
- Even a **tiny change** in input creates completely different output
- Is efficient to compute

SHA256(hello) =
2cf24dba5fb0a30e26e83b2ac5
b9e29e1b161e5c1fa7425e7304
3362938b9824
SHA256(hullo) =
7835066a1457504217688c8f5d
06909c6591e0ca78c254ccf174
50d0d999cab0

**Note:** One character change  $\rightarrow$  completely different hash!

## Expected properties of a hash function

- **Collision resistance**: computationally infeasible to find two different inputs producing the same hash.
- **Preimage resistance**: given the output of a hash function, it is computationally infeasible to reconstruct the original input.
- Second preimage resistance: given an input and an output, it's computationally infeasible to find another different input producing the same output.



SHA-2 compression function. Source: Wikipedia

### Hash functions: what are they good for?

- **Password storage**: Store the hash of the password on the server, not the password itself. Then check candidate passwords against the hash.
- **Data integrity verification**: Hash a file. Later hash it again and compare hashes to check if the file has changed, suffered storage degradation, etc.
- **Proof of work**: Server asks client to hash something a lot of times before they can access some resource. Useful for anti-spam, Bitcoin mining, etc.
- Zero knowledge proofs: time for more actual magic

## Time for more actual magic

- Zero-knowledge proofs allow you to prove that you know a secret without revealing any information about it.
- They built "zero-knowledge virtual machines" where you can execute an entire program that runs as a zero-knowledge proof.
- ZKP battleship game: server proves to the players that its output to their battleship guesses is correct, without revealing any additional information (e.g. ship location).



Battleship board game. Source: Hasbro

### Evaluating a hash function's quality

- Recall:
  - Asymmetric primitives are based on mathematical problems, can be mathematically proven secure (given assumptions!)
  - Symmetric primitives (encryption, hashing...) are statistically, empirically, heuristically shown to be secure, not proven secure.
  - The more cryptanalysis they survive, the higher confidence we have in their security.



Cryptanalysis of AES.

### What about encryption?

- Symmetric primitive of choice for encryption: **AES**.
- Not that far off in terms of design process from hash functions, but:
  - AES is a PRP (pseudorandom permutation)
  - HMAC-SHA256 is a PRF (pseudorandom function)



AES's SubBytes operation. Source: Wikipedia

#### **PRF versus PRP**

#### Pseudo-Random Function (SHA-2)

- Input is arbitrary-length,
- **Output** is fixed-length, looks random (as discussed earlier).
- Indistinguishable from a truly random function by an adversary with limited computational power.

#### Pseudo-Random Permutation (AES)

- Input and output are the same length, forming a permutation.
- Each input maps uniquely to one output, allowing invertibility.
- Indistinguishable from a truly random permutation by an adversary with limited computational power.

 $\mathsf{PRF} : F_k = X \to Y$ 

- We want the mapping to be:
  - One-way
  - "Randomized"
  - Relations between inputs not reflected in outputs



 $\mathsf{PRP} : F_k = X \to X$ 

- Bijective (two-way)
  - Injective: no two inputs map to same output (no collisions)
  - Surjective: Every output has one corresponding input
- "Randomized"
- Relations between inputs not reflected in outputs



#### AES is a block cipher

- AES takes a 16-byte input, produces a 16-byte output.
- Key can be 16, 24 or 32 bytes.
- OK, so what if we want to encrypt more than 16 bytes?
- **Proposal**: split the plaintext into 16 byte chunks, encrypt each of them with the same key.

### Block cipher examples







What we actually want

What we start with

What we get

## Block cipher modes of operation



Source: Wikipedia

# Cryptographic building blocks

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## TLS 1.3: high-level sketch



#### Source: Mostafa Ibrahim

### TLS 1.3: high-level sketch

- **Public key agreement** (eg. Diffie-Hellman) is used to establish a shared secret between the client and the server.
- **AES** is used for encrypting data in transit.
- SHA-2 is used for hashing (checking certificates, etc.)





### TLS 1.3: high-level sketch

- Through the design, we accomplish the desired security goals under a well-specified threat model:
- Security goals: confidentiality of data, authentication of the server towards the client...
- Threat model: malicious Internet Service Provider (ISP), etc.





#### How TLS 1.3 was made



#### How TLS 1.3 was made



### From hard problems to real-world security

#### The journey we'll trace

- 1. Mathematical insight: Discrete logarithm is hard to compute.
- 2. Cryptographic innovation: Diffie-Hellman key exchange leverages this hardness.
- 3. Real-world impact: Secure communication for billions of people daily.

**This is the power of applied cryptography**: transforming abstract mathematical problems into tools that help people and protect our digital lives.

#### Course goals

- Understand the reasoning behind the math of modern cryptography.
- Analyze and prove the security of cryptographic constructions.
- Understand how cryptographic constructions can be composed to build real-world secure protocols and systems.
- Discern between theoretical cryptography and applied cryptography from an engineering perspective.
- Critically assess security implementations and evaluate real-world cryptographic protocols.
- Gain an understanding of the future of cryptography and its role in emerging technologies.

#### **Course prerequisites**

- Good but optional: CMPS 215 (Theory of Computation)
- If you want to understand whether you have the sufficient background for this course, review this revision chapter and try to do all the exercises: https://joyofcryptography.com/pdf/chap0.pdf

#### **Class materials**

- Joy of Cryptography: learn how to reason about and prove systems secure.
- Attack papers, codebases, labs: hard engineering perspective.

- Always keep an eye on the website: Course news, updates, materials, slides will all be posted there. https://appliedcryptography.page
- I am aiming for the most engaging course possible!







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